

# Estimating Cardinality: Use of Jonathan Lewis CBO methodology

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#### **Cost-Based Oracle Fundamentals**

By Jonathan Lewis ISBN13: 978-1-59059-636-4 ISBN10: 1-59059-636-6 536 pp. Published Oct 2005 by <u>Apress</u>

# Cost-Based Oracle Fundamentals

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"The insights that Jonathan provides into the workings of the cost-based optimizer will make a DBA a better designer and a developer a better SQL coder. Both groups will become better troubleshooters." —Thomas Kyte

#### **Jonathan Lewis**

Foreword by Thomas Kyte Vice President (Public Sector), Oracle Corporation



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#### How many red M&Ms?



Estimate total count Count colors Assume even distribution Divide total by colors

> (NoCOUG <u>Physical Data</u> <u>Storage presentation on</u> November 11, 2009)



# **Cardinality and Selectivity**

- Selectivity = fraction of all rows matched by a predicate (or predicates)
- Cardinality = Selectivity \* Rows



## **ClientX Group Members Query**

```
-- 61zyzhy11juwf
select
   cons id,
   member id,
   active,
   first name,
   last name,
   primary email,
   cat id,
   admin status
from
   clientX.constituent cons
   inner join clientX.group members gm
      on gm.user id = cons.cons id
     and gm.group_id = :"SYS_B_0"
where active = :"SYS B 1"
and cons.site id = :"SYS B 2"
and rownum <= :"SYS B 3"
ORDER BY
   last name lower ASC,
   first name lower ASC;
```



#### Plan hash value: 990099801

Id	Operation	Name	Rows	Bytes
0     1    * 2    * 3    * 4    * 5	SELECT STATEMENT   SORT ORDER BY   COUNT STOPKEY   HASH JOIN   INDEX RANGE SCAN   TABLE ACCESS FULL	GROUP_USER_INDEX CONSTITUENT	   2001   144K   144K   144K   4450K	1552K

Peeked Binds (identified by position):

1 - :SYS\_B\_0 (NUMBER): 77569 2 - :SYS\_B\_1 (NUMBER): 1 3 - :SYS\_B\_2 (NUMBER): 1638 4 - :SYS\_B\_3 (NUMBER): 2001

Predicate Information (identified by operation id):

```
2 - filter(ROWNUM<=:SYS_B_3)
3 - access("GM"."USER_ID"="CONS_ID")
4 - access("GM"."GROUP_ID"=:SYS_B_0)
5 - filter(("ACTIVE"=:SYS_B_1 AND "SITE_ID"=:SYS_B_2))</pre>
```



#### **Plan narrative**

- Scans GROUP\_USER\_INDEX for the GROUP\_ID specified by bind variable.
- Reads index blocks, and constructs a hash table of USER\_ID values that are in this group.
- Full scans CONSTITUENT, hashing each CONS\_ID, feeding matching rows up to parent.
- Stops full scan after 2001 matches.
- Sorts the 2001 rows, passes them to Java



# How many constituents in a group?

- Look at "Rows" column for Operation #4 Oracle estimates 144,000 rows for access("GM"."GROUP\_ID"=:SYS\_B\_0)
- Jonathan Lewis describes two methods (p. 43):
   1. cardinality = (rows in table) / (distinct values)
   2. cardinality = (rows in table) \* (density)
- Which method gets used is not documented by Oracle, and can vary with version, etc.
- Method #2 is usually used with histograms, which do exist on this column



# How many constituents in a group?

Value	Description
0.159%	DBA_TAB_COLUMNS.DENSITY of clientX.group_members.group_id
90,641,999	DBA_TABLES.NUM_ROWS (number of rows) in clientX.group_members
725	DBA_TAB_COLUMNS.NUM_DISTINCT (number of distinct values) in clientX.group_members.group_id
125,023	estimated cardinality of arbitrary group_id, <b>method 1, rows/distinct</b> (p. 43)
144,565	estimated cardinality of arbitrary group_id, <b>method 2, density * rows</b> (p. 43)



#### Method 2 was used

- Oracle used Method 2 (rows \* density)
   144,565 equals the 144K in the execution plan
- In other words, the mean number of constituents in a group is about 144,000 (or 125,000, depending on estimation method).
- Two methods are usually equivalent.
- Watch out when DENSITY \* NUM\_DISTINCT is far from unity (here 0.159% \* 725 = 1.15, OK)



## Best way to get CONS data for 144K rows

- Oracle estimates 144K values of USER\_ID.
- Need to find CONSTITUENT rows with matching CONS\_ID, and that satisfy other predicates
- Two main choices:

Full scan and do a hash join, or Nested loop into an index with table access by rowid



## Full scan: multi-block

- Full scans are done with multi-block reads
- A multi-block read is more expensive than a single block read, but
- The cost per block is less when done via multiblock than when done one at a time.
- Based on system performance stats stored in sys.aux\_stats\$ (to which I had no access)



# Multi-block read efficiency (Lewis, p. 20)

Metric	Value	Description (typical or actual)
IOSEEKTIN	10	Disk seek time in milliseconds (typical)
IOTFRSPEED	4096	Disk transfer time in bytes per millisecond (typical)
DB_BLOCK_SIZE	8192	Block size in bytes (actual)
DB_FILE_MULTIBLOCK_READ_COUNT	16	How many blocks are read in a multi- block read (actual)
SREADTIM	(tbd)	Single block read time in milliseconds (to be calculated)
MREADTIM	(tbd)	Multi-block read time (total) in milliseconds (to be calculated)

sreadtim = ioseektim + db\_block\_size/iotrfrspeed

mreadtim = ioseektim

+ db\_file\_multiblock\_read\_count \*(db\_block\_size/iotrfrspeed)



# Multi-block read efficiency (Lewis p. 20-1)

Value	Description	
10	ioseektim (typical)	
8192	db_block_size (actual)	
4096	iotrfrspeed (typical)	
16	db_file_multiblock_read_count (actual)	
12	<pre>sreadtim = ioseektim + db_block_size/iotrfrspeed (Lewis, pp. 20)</pre>	
42	<pre>mreadtim = ioseektim + (db_file_multiblock_read_count *(db_block_size/iotrfrspeed))</pre>	
3.5	mreadtim/sreadtim	



# Full scan cost: clientX.constituent

Value	Description	
608,252	number of blocks in clientX.CONSTITUENT	
16	db_file_multiblock_read_count (actual)	
38,016	multi-block-reads, number of IO requests to do full scan: (number of blocks / blocks per read)	
3.5	mreadtim/sreadtim	
133,055	equivalent number of single-block IO requests for full table scan access: multi-block-reads * (mreadtime/sreadtim)	

# A full scan costs 608,252 "gets", but costs only as much as 133,055 single-block gets



#### **Nested loop cost**

- Single-block reads
- Depends on:

Availability of suitable CONS\_ID index Number of rows per CONS\_ID Cost of reading index (per CONS\_ID) Cost of reading table blocks (per CONS\_ID) Number of CONS\_ID values to lookup



# Index suitability

- No primary key **index** here!
- Only index leading with CONS\_ID: ITOPS\_CONS\_ID\_UNAME
- But this index is not unique how many rows per CONS\_ID?

 Oracle can estimate rows/CONS\_ID by: Table does have a CONS\_ID primary key Table stats (explained below) Index stats (explained below)



#### Index suitability – table stats

From dba\_tab\_columns

DENSITY of CONS\_ID = 2.2151E-07

NUM\_DISTINCT of CONS\_ID = 4,514,534

NUM\_ROWS of CONSTITUENT = 4,532,476

computed cardinality method 1 (no histogram), num\_rows/num\_distinct, p 43 = **1.00397** 

computed cardinality method 2 (if histogram), num\_rows \* density, p 43 = **1.00399** 



## Index suitability – index stats

From dba\_indexes:

DISTINCT\_KEYS of ITOPS\_CONS\_ID\_UNAME = 4,761,265 NUM\_ROWS of ITOPS\_CONS\_ID\_UNAME = 4,761,265 computed cardinality = NUM\_ROWS/DISTINCT\_KEYS = 1

 All three methods indicate that each CONS\_ID has a single row, so a NESTED LOOP seems potentially feasible.



# Index/table costs per CONS\_ID

- BLEVEL clientX.ITOPS\_CONS\_ID\_UNAME = 2
- 2 reads to get to leaf block
- 1 read to get leaf block
- 1 read to get table block (from index rowid)
- 4 reads per CONS\_ID



# **Nested loop cost: clientX.constituent**

Value	Description
144,565	Estimated CONS_ID values to lookup (from previous discussion)
4	Single-block reads per CONS_ID
578,258	Single-block reads to do all lookups

Previously shown: a full scan costs **608,252** "gets", but costs only as much as **133,055** single-block gets.

But a nested loop costs **578,258** single-block reads, so is much more expensive.

Note that this analysis <u>depends heavily on the number of</u> <u>CONS\_ID values to lookup</u>.

# What about logical vs. physical?

- Optimizer ignores the difference between
  - blocks that are cached in the SGA, and
  - those that must be read from disk.
- Hard to account for this difference in a simple calculation, due to the variable nature of the cache. Assumes all reads are physical.
- However, this assumption is probably pretty good for large table scans, whose blocks are least likely to remain in the cache.



#### Just scratched the surface

- Study: range scans and clustering factor
- Study: ranges and high/low value
- Take home: estimated row counts are key!
- Cardinality Feedback



#### How many red M&Ms?



Estimate total count Count colors Assume even distribution Divide total by colors **Physical Data Storage** November 11, 2009 SQL Execution Plans, DBMS\_XPLAN, and Cardinality Feedback May 20, 2010

